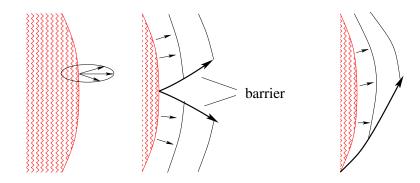
# Dynamic Blocking Problems for a Model of Fire Propagation

#### Alberto Bressan

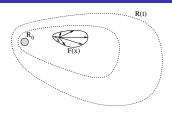
Mathematics Department, Penn State University

(Waterloo, 2011)

# Blocking an advancing wildfire



## A Differential Inclusion Model for Fire Propagation



$$R(t) \subset \mathbb{R}^2$$
 = set reached by the fire at time  $t \geq 0$ 

determined as the reachable set by a differential inclusion

$$\dot{x} \in F(x)$$
  $x(0) \in R_0 \subset \mathbb{R}^2$ 

Fire may spread in different directions with different velocities

$$R(t) = \begin{cases} x(t); & x(\cdot) \text{ absolutely continuous}, \end{cases}$$

$$x(0) \in R_0$$
,  $\dot{x}(\tau) \in F\big(x(\tau)\big)$  for a.e.  $\tau \in [0,t]$ 



## Confinement Strategies

(A.B., J.Differential Equations, 2007)

Assume: a **controller** can construct a **wall**, i.e. a one-dimensional rectifiable curve  $\gamma$ , which blocks the spreading of the fire.

 $\gamma(t)\subset\mathbb{R}^2=$  portion of the wall constructed within time t  $\sigma=$  speed at which the wall is constructed

**Definition 1.** A set valued map  $t\mapsto \gamma(t)\subset \mathbb{R}^2$  is an admissible strategy if :

(H1) For every  $t_1 \leq t_2$  one has  $\gamma(t_1) \subseteq \gamma(t_2)$ 

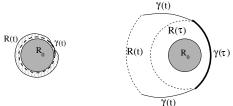
(H2) Each  $\gamma(t)$  is a rectifiable set (possibly not connected). Its length satisfies

$$m_1(\gamma(t)) \leq \sigma t$$

**Definition 2.** The reachable set determined by the blocking strategy  $\gamma$  is

$$R^{\gamma}(t) \doteq \left\{ x(t); \ x(\cdot) \ \text{ absolutely continuous}, \ x(0) \in R_0 
ight.$$
  $\dot{x}( au) \in Fig(x( au)ig) \ \text{ for a.e. } au \in [0,t], \qquad x( au) 
otin \gamma(t) \quad \text{for all } au \in [0,t] \ 
ight\}$ 

REMARK: Walls must be constructed in real time!



An admissible strategy is described by a set-valued function  $t\mapsto \gamma(t)\subset \mathbb{R}^2$ 

 $\gamma(t)$  = portion of the wall constructed within time t

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# Static vs. Dynamic Blocking Problems

## Static problems



minimizing enclosed area



minimizing length



minimizing area + length

#### **Dynamic problem**









## **Optimal Confinement Strategies**

#### A cost functional should take into account

- The value of the region destroyed by the fire.
- The cost of building the wall.
- $\alpha(x)$  = value of a unit area of land around the point x
- $\beta(x) = \cos t$  of building a unit length of wall near the point x

#### COST FUNCTIONAL

$$J(\gamma) \doteq \lim_{t \to \infty} \left\{ \int_{R^{\gamma}(t)} \alpha \, dm_2 + \int_{\gamma(t)} \beta \, dm_1 \right\}$$

### Mathematical Problems

### 1. Blocking Problem.

Given an initial set  $R_0$ , a multifunction F and a wall construction speed  $\sigma$ , does there exist an admissible strategy  $t\mapsto \gamma(t)$  such that the reachable sets  $R^{\gamma}(t)$  remain uniformly bounded for all t>0?

#### 2. Optimization Problem.

Find an admissible strategy  $\gamma \in \mathcal{S}$  which minimizes the cost functional  $J(\gamma)$ .

- (i) Existence of an optimal solution
- (ii) **Necessary conditions** for the optimality of an admissible strategy  $\gamma(\cdot)$
- (iii) **Regularity** of the curves  $\gamma(t)$  constructed by an optimal strategy
- (iv) Sufficient conditions for the optimality of a strategy  $\gamma(\cdot)$
- (v) Numerical computation of an optimal strategy

## Equivalent Formulation

(A.B. - T. Wang, Control Optim. Calc. Var. 2009)

Blocking Problems and Optimization Problems can be reformulated in terms of one single rectifiable set  $\Gamma$ 

$$(\mathsf{strategy}) \qquad t \mapsto \gamma(t) \qquad \longleftrightarrow \qquad \mathsf{\Gamma} \qquad (\mathsf{single \ wall})$$

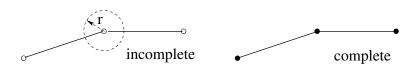
$$\gamma(\cdot) \longrightarrow \Gamma \doteq \left(\bigcup_{t>0} \gamma(t)\right) \setminus \left(\bigcup_{t>0} R^{\gamma}(t)\right) \quad \text{(useful walls)}$$

$$\Gamma \longrightarrow \gamma(t) \doteq \Gamma \cap \overline{R^{\Gamma}(t)}$$
 (walls touched by the fire within time  $t$ )

## Complete strategies

A rectifiable set  $\Gamma$  is **complete** if it contains all of its points of positive upper density:

$$\theta^*(\Gamma;x) \; \doteq \; \limsup_{r \to 0+} \frac{m_1\Big(B(x,r) \cap \Gamma\Big)}{r} \; > \; 0 \qquad \Longrightarrow \qquad x \in \Gamma \, .$$



The **reachable sets** for the differential inclusion

$$\dot{x} \in F(x)$$
  $x(0) \in R_0$   $(DI)$ 

without crossing  $\Gamma$  are

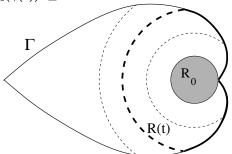
$$R^{\Gamma}(t) \doteq \left\{ x(t); \ x(\cdot) \text{ absolutely continuous}, \ x(0) \in R_0 \right.$$
  $\dot{x}(\tau) \in F(x(\tau)) \text{ for a.e. } \tau, \qquad x(\tau) \notin \Gamma \text{ for all } \tau \in [0,t] \right\}$ 

#### Admissible curves

The complete, rectifiable set  $\Gamma$  is **admissible** for the differential inclusion (DI) and the construction speed  $\sigma$  if, for every  $t \ge 0$ , the set

$$\gamma(t) \doteq \Gamma \cap \overline{R^{\Gamma}(t)}$$

has total length  $m_1(\gamma(t)) \leq \sigma t$ 



 $\gamma(t)$  = portion of  $\Gamma$  that needs to be in place by time t

## **Equivalent Formulations**

$$R_{\infty}^{\Gamma} \doteq \bigcup_{t>0} R^{\Gamma}(t) = [\text{total region burned by the fire}]$$

**Blocking Problem 2:** Find an admissible rectifiable set  $\Gamma \subset \mathbb{R}^2$  such that the corresponding reachable set  $R^{\Gamma}_{\infty}$  is bounded.

**Optimization Problem 2:** Find an admissible rectifiable set  $\Gamma \subset \mathbb{R}^2$  which minimizes the cost

$$J(\Gamma) = \int_{R_{\infty}^{\Gamma}} \alpha \, dm_2 + \int_{\Gamma} \beta \, dm_1$$

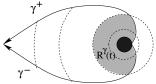
## Blocking the Fire

- Fire propagates in all directions with unit speed:  $F(x) = B_1$
- ullet Wall is constructed at speed  $\sigma$

## **Theorem** (A.B., J.Differential Equations, 2007)

On the entire plane, the fire can be blocked if  $\sigma > 2$ , it cannot be blocked if  $\sigma < 1$ .

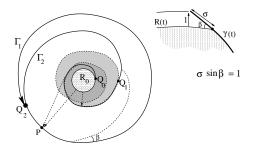
**Blocking Strategy:** If  $\sigma > 2$ , construct two arcs of logarithmic spirals along the edge of the fire



$$\gamma(t) \doteq \left\{ (r, \theta); \quad r = \mathrm{e}^{\lambda |\theta|}, \quad 1 \leq r \leq 1 + t 
ight\}, \qquad \lambda \doteq \frac{1}{\sqrt{rac{\sigma^2}{4} - 1}}$$

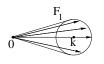
### When can the fire be blocked?

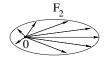
**Conjecture:** Assume the fire propagates with speed 1 in all directions. On the entire plane the fire can be blocked if and only if  $\sigma > 2$ 

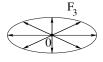


Single spiral strategy: curve closes on itself if and only if  $\sigma > \sigma^{\dagger} = 2.614430844\dots$  (M. Burago, 2006)

## Non-isotropic fire propagation







**Theorem.** (A.B., M. Burago, A. Friend, J. Jou, Analysis and Applications, 2008) If the wall construction speed satisfies

$$\sigma \ > \ [ \text{vertical width of} \ F ] \ = \ 2 \max_{\theta \in [0,\pi]} \ \rho(\theta) \sin \theta$$

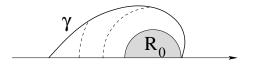
then, for every bounded initial set  $R_0$ , a blocking strategy exists

## The isotropic case on the half plane

- Fire propagates in all directions with unit speed.  $F(x) = B_1$
- ullet Wall is constructed at speed  $\sigma$

## Theorem. (A.B. - T.Wang, J.Math Anal.Appl. 2009)

Restricted to a half plane, the fire can be blocked if and only if  $\sigma>1$ 

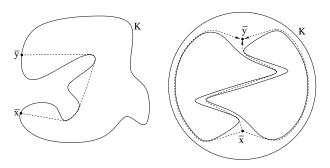


BLOCKING STRATEGY: If  $\sigma>1$ , the fire can be enclosed between the horizontal axis and an arc of logarithmic spiral

$$\lambda \doteq \frac{1}{\sqrt{\sigma^2 - 1}}$$
  $\Gamma \doteq \left\{ (r, \theta); \quad r = e^{\lambda \theta}, \quad \theta \in [0, \pi] \right\}$ 

## Regularity of the distance function

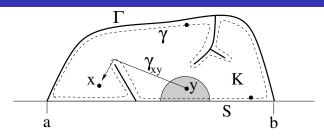
 $d_K(x,y) \doteq \text{minimum length among all paths } \gamma : [0,1] \mapsto K \text{ joining } x \text{ with } y$ 



**Lemma.** If  $K \subset \mathbb{R}^2$  is compact, simply connected, then the map  $y \mapsto d_K(x,y)$  is  $C^1$  in the interior of  $K \setminus \{x\}$ .

 $\sup_{x,y\in K} d_K(x,y)$  is attained at boundary points.

### Estimates on the distance function



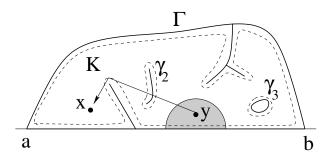
Assume: barrier  $\Gamma$  is completed at time T > 0

S =segment joining the points a, b

Any two points  $x,y\in K$  can be connected by a path  $\gamma_{xy}$  of length

$$m_1(\gamma_{xy}) \leq \frac{1}{2}m_1\bigg(\Gamma \cup S\bigg) \leq \frac{1}{2}\bigg(\sigma T + (b-a)\bigg) < \sigma T$$

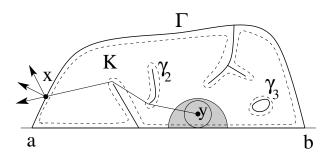
Same conclusion in the case of several connected components



Any two points  $x, y \in K$  can be connected by a path  $\gamma_{xy}$  of length

$$m_1(\gamma_{\mathsf{x}\mathsf{y}}) \ \le \ \frac{1}{2} m_1 igg( \Gamma \cup \mathcal{S} igg) \ \le \ \frac{1}{2} igg( \sigma \, T + (b-a) igg) \ < \ \sigma \, T$$

# No strategy can block the fire if $\sigma \leq 1$



x = position of "last brick of the wall"

Fire reaches x and spreads outside, before time T when the barrier is completed

$$d_K(x,y) \leq \max_{x \in \partial K} d_K(x,y) \leq \frac{1}{2} m_1(\Gamma \cup S) \leq m_1(\Gamma) \leq \sigma T$$

## **Existence of Optimal Strategies**

Fire propagation:  $\dot{x} \in F(x)$   $x(0) \in R_0$ 

Wall constraint:  $\int_{\gamma(t)} \psi \, dm_1 \leq t$   $(1/\psi(x) = \text{construction speed at } x)$ 

Minimize:  $J(\gamma) = \left\{ \int_{R^{\gamma}(t)} \alpha \, dm_2 + \int_{\gamma(t)} \beta \, dm_1 \right\}$ 

#### **Assumptions:**

- (A1) The initial set  $R_0$  is open and bounded. Its boundary satisfies  $m_2\left(\partial R_0\right)=0$ .
- (A2) The multifunction F is Lipschitz continuous w.r.t. the Hausdorff distance. For each  $x \in \mathbb{R}^2$  the set F(x) is compact, convex, and contains a ball of radius  $\rho_0 > 0$  centered at the origin.
- (A3) For every  $x \in \mathbb{R}^2$  one has  $\alpha(x) \ge 0$ ,  $\beta(x) \ge 0$ , and  $\psi(x) \ge \psi_0 > 0$ .  $\alpha$  is locally integrable, while  $\beta$  and  $\psi$  are both lower semicontinuous.

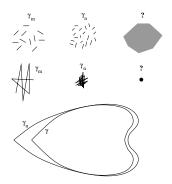
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#### **Theorem** (A.B. - C. De Lellis, *Comm. Pure Appl. Math.* 2008)

Assume (A1)-(A3), and  $\inf_{\gamma \in \mathcal{S}} J(\gamma) < \infty$ .

Then the minimization problem admits an optimal solution  $\gamma^*$ .

**Direct method:** Consider a minimizing sequence of strategies  $\gamma_n(\cdot)$  Define the optimal strategy  $\gamma^*$  as a suitable limit.



**STEP 1:** Replace the  $\gamma_n(\cdot)$  by **complete strategies**  $\gamma_n^c$ 

$$\theta^*(E,x) \doteq \limsup_{r\downarrow 0} \frac{m_1(B(x,r)\cap E)}{2r}$$
 (upper density)

$$\overline{\gamma}_n(t) \ \doteq \ \gamma_n(t) \ \cup \ \left\{ x \in \mathbb{R}^2 \, ; \quad \theta^*(\gamma_n(t), x) > 0 \right\}, \qquad \qquad \gamma_n^c(t) \ \doteq \ \bigcap_{s > t} \overline{\gamma}_n(s)$$

**STEP 2:** For each rational time  $\tau$ , order the connected components of  $\gamma_n(\tau)$ according to decreasing length:

$$\ell_{n,1} \ge \ell_{n,2} \ge \ell_{n,3} \ge \cdots$$
  $\gamma_n(t) = \gamma_{n,1} \cup \gamma_{n,2} \cup \gamma_{n,3} \cup \cdots$ 

Taking a subsequence, as  $n \to \infty$  we can assume

$$\ell_{n,i}( au) 
ightarrow \ell_i( au) \qquad \qquad \gamma_{n,i}( au) 
ightarrow \gamma_i( au) \qquad \qquad au \in \mathbf{Q}$$

We then define  $\gamma(\tau) \doteq \bigcup \gamma_i(\tau)$ . Finally  $\gamma^*(\cdot) \doteq$  completion of  $\gamma(\cdot)$ .  $i > 1, \ell_i(\tau) > 0$ 

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**STEP 3:** Lower semicontinuity of  $\psi, \beta$  imply

$$\int_{\gamma^*(t)} \psi \, dm_1 \, \leq \, \liminf_{n \to \infty} \int_{\gamma_n(t)} \psi \, dm_1 \, \leq \, t \qquad \qquad \int_{\gamma^*(t)} \beta \, dm_1 \leq \liminf_{n \to \infty} \int_{\gamma_n(t)} \beta \, dm_1$$

Hence the limit strategy  $\gamma(\cdot)$  is admissible.

#### STEP 4:

$$\int_{R^{\gamma^*}(t)} \alpha \, dm_2 \, \leq \, \liminf_{n \to \infty} \int_{R^{\gamma_n}(t)} \alpha \, dm_2$$

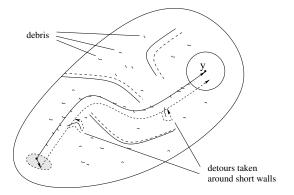
Claim: given  $y \in R^{\gamma^*}(t)$  and r > 0, for every n large:  $R^{\gamma_n}(t) \cap B(y,r) \neq \emptyset$ 

Given 
$$\varepsilon > 0$$
, decompose:  $\gamma_n(\tau) = \left(\bigcup_{i \leq i_1(\tau)} \gamma_{n,i}^{\tau}\right) \cup \gamma_n'(\tau) \cup \gamma_n''(\tau)$ 

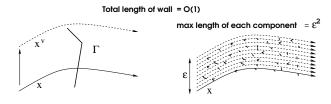
 $[large\ connected\ components]\ \cup\ [small\ connected\ components]\ \cup\ [debris]$ 

$$m_1(\gamma_n'(\tau)) < \varepsilon$$
  $m_1(\gamma_n''(\tau)) = \mathcal{O}(1)$ 

 $\lim_{n\to\infty} [\text{maximum length of connected components in } \gamma_n''] = 0$ 

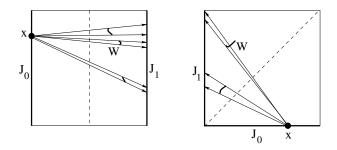


- Shift the trajectory  $x^{\sharp}$  so that it crosses an amount  $\mathcal{O}(\varepsilon)$  of debris (connected components of  $\gamma_n''$ ) for all n.
- Take detours of total length  $\mathcal{O}(\varepsilon)$  to get around the small connected components  $\gamma_n'$  and the debris.



Given  $x(\cdot)$ , find a shifted trajectory  $x^{\nu}(\cdot)$  so that  $[\text{total length of connected components of }\Gamma\text{ crossed by }x^{\nu}(\cdot)]\ <\ \varepsilon$ 

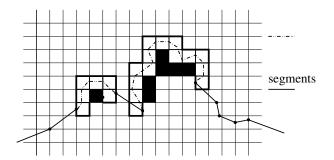
**Lemma 1:** In case 2, this can be achieved by a shift  $|v| = \mathcal{O}(\varepsilon)$ 



**Lemma 2:**  $J_0, J_1$  edges of the unit square Q. For any  $\mu > 0$  there exists  $\kappa > 0$  such that:

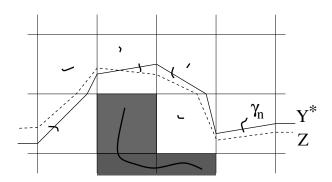
If  $W \subset Q$  is a set with  $m_1(W) < \kappa$ , then there is a set  $K_0 \subseteq J_0$ 

- (i)  $m_1(K_0) \geq 1 \mu$ ,
- (ii) For every  $x \in K_0$ , the set  $J_x$  of points  $y \in J_1$  for which the segment [x, y] with endpoints x, y does not intersect W has measure  $m_1(J_x) \ge 1 \mu$ .



- Construct a  $\sigma$ -grid
- ullet Replace trajectory by a polygonal  $Y^\sigma$
- If  $m_1(\gamma_n \cup Q) > \kappa \sigma$  the square Q is black. Otherwise it is white.
- ullet Construct polygonal detours around the black islands, with total length  $\mathcal{O}(arepsilon)$

ullet Modify the trajectory  $Y^*$  inside white squares so that the new polygonal Z does not cross any wall of  $\gamma_n$ 



## Alternative approach: the minimum time function

(C.DeLellis & R.Robyr, 2010)

$$\dot{x} \in F(x)$$
  $x(0) \in R_0 \subset \mathbb{R}^2$ 

Assume:  $B(0, \rho) \subseteq F(x)$ 

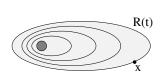
Minimum time function: 
$$T(x) \doteq \inf\{t \geq 0; x \in R(t)\}$$

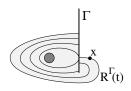
is Lipschitz continuous and satisfies the Hamilton-Jacobi equation

$$H(x, \nabla T(x)) = 0$$
 a.e.

$$H(x, p) = \max_{v \in F(x)} \langle v, p \rangle - 1$$

## The minimum time function with obstacle





**Barrier:** a complete, rectifiable set  $\Gamma \subset \mathbb{R}^2$ 

 $T^{\Gamma}(x) \doteq \inf\{t > 0 : x \in R^{\Gamma}(t)\}$ Minimum time function with obstacle:

 $R^{\Gamma}(t) = \text{set of points reached by } F\text{-trajectories which start in } R_0$ and do not cross [

**Goal:** characterize  $T^{\Gamma}$  as the unique maximal subsolution to a H-J equation

# A family $S^{\Gamma}$ of subsolutions

**Definition.** A function  $u: \mathbb{R}^2 \mapsto [0, \infty]$  is in the set  $\mathcal{S}^{\Gamma}$  if

- $u \in SBV$   $(Du = \nabla u + D^{jump}u + D^{Cantor}u, \text{ with } D^{Cantor}u \equiv 0)$
- $m_1(J_u \setminus \Gamma) = 0$   $(J_u \doteq \text{set of jump points of } u)$
- u = 0 on  $R_0$
- $H(x, \nabla u(x)) \leq 0$  for a.e. x

 $abla u \doteq {\sf absolutely}$  continuous part of the distributional derivative  ${\sf D} u$ .

## Theorem. (C.DeLellis & R.Robyr, 2010 - Archive Rat. Mech. Anal.)

Let  $\Gamma$  be a complete, rectifiable set. Then the minimum time function  $\mathcal{T}^{\Gamma}$  is the unique maximal element of  $\mathcal{S}^{\Gamma}$ .

## Corollary. (C.DeLellis & R.Robyr, 2010)

Consider the cost functions  $\alpha \geq 0, \beta \geq 0$ , with  $\alpha$  locally integrable and  $\beta$  lower semicontinuous. Then there exist a blocking strategy  $\Gamma$  which minimizes the cost among all admissible ones.

**Proof of the Corollary.** Take a minimizing sequence of admissible barriers  $\Gamma_k$ .

Consider the corresponding minimum time functions  $T_{\iota} \doteq T^{\Gamma_{\iota}}$ 

Using the Ambrosio-De Giorgi compactness theorem for SBV functions, one obtains a convergent subsequence  $T_k \to U$ , such that

- $U \in SBV$
- the jump set  $J_U$  is a rectifiable set, and

$$\int_{J_U} \psi \, dm_1 \, \leq \, \liminf_{k \to \infty} \, \int_{J_{T_k}} \psi \, dm_1 \, \leq \, t \, .$$

Hence  $\Gamma \doteq [\text{completion of } J_U]$  is an admissible barrier.

$$\int_{R_{\infty}^{\Gamma}} \alpha \ dm_2 \ \leq \ \liminf_{k \to \infty} \int_{R_{\infty}^{\Gamma_k}} \alpha \ dm_2 \qquad \qquad \int_{\Gamma} \alpha \ dm_2 \ \leq \ \liminf_{k \to \infty} \int_{\Gamma_k} \alpha \ dm_2$$

## Necessary conditions for optimality

**Problem:** find an admissible barrier  $\Gamma$  which minimizes

$$J(\Gamma) \doteq \alpha \cdot [\text{total burned area}] + \beta \cdot [\text{length of the curve}]$$

**GOAL:** derive a set of ODE's describing the walls built by an optimal strategy

- A.B., J.Differential Equations, 2007
- A.B. T.Wang, ESAIM, Control Optim. Calc. Var. 2010
- T.Wang, *Intern. J. Control*, to appear.

## Classification of arcs in an optimal strategy

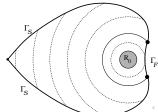
Minimum time function  $T^{\Gamma}(x) \doteq \inf \left\{ t \geq 0 \; ; \; x \in \overline{R^{\Gamma}(t)} \right\}$ 

Set of times where the constraint is saturated

$$\mathcal{S} \doteq \left\{ t \geq 0 \, ; \; \; \mathit{meas} \left( \Gamma \cap \overline{R^{\Gamma}(t)} \right) \; = \; \sigma t \, \right\}$$

**Boundary arcs:**  $\Gamma_{\mathcal{S}} \doteq \{x \in \Gamma; \quad \mathcal{T}^{\Gamma}(x) \in \mathcal{S}\}$  constructed along the advancing fire front

Free arcs:  $\Gamma_{\mathcal{F}} \doteq \{x \in \Gamma; T^{\Gamma}(x) \notin \mathcal{S}\}$  constructed away from the fire front



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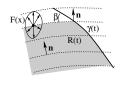
# Optimality conditions, minimizing the value of burned area

1. A free arc  $\ \Gamma$ . The curvature must be proportional to the local value of the land

$$r(s) = \text{radius of curvature} \quad \alpha = \text{land value}$$

$$r(s) \cdot \alpha(\Gamma(s)) = \text{const.}$$

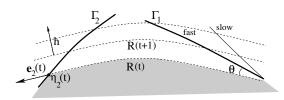




2. A single boundary arc  $\Gamma$ . The wall is constructed at maximum speed  $\sigma$ , always remaining at the edge of the burned set

$$\sigma \sin \beta = \max_{y \in F(x)} \mathbf{n} \cdot y$$

# **3. Two or more boundary arcs:** $\Gamma_1, \ldots, \Gamma_{\nu}$ , constructed simultaneously for $t \in [a, b]$



Sum of construction speeds  $< \sigma$ 

At which speed should each wall be constructed ?

(Solution is found applying Pontryagin's Maximum Principle)

#### Value of Time

$$J(\Gamma) \doteq \alpha \cdot [\text{total burned area}] + \beta \cdot [\text{length of the curve}]$$

There exists a non-increasing scalar function  $t \mapsto W(t)$ 

 $\approx$  Lagrange multiplier corresponding to the constraint

$$m_1 \Biggl( \Gamma \cap \overline{R^\Gamma(t)} \Biggr) \ \le \ \sigma t$$

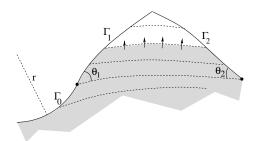
1. A free arc  $\Gamma_0$ :  $W(t) = (\alpha r - \beta)\sigma = \text{constant}$ 

r = radius of curvature

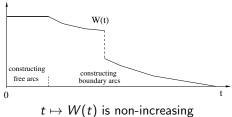
#### 2. Two or more boundary arcs: $\Gamma_1, \ldots, \Gamma_{\nu}$ , constructed simultaneously

$$W(t) = W_i(t) = \left(\frac{q_i(t)}{\cos\theta_i(t)} - \beta\right)\sigma$$

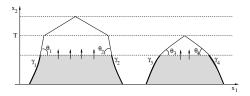
 $q = (q_1, q_2, \dots, q_{\nu}) =$  adjoint variable in the PMP.



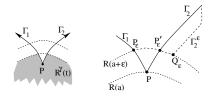
## Instantaneous value of time



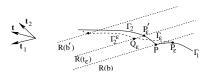
- (')
- continuous at times where a free arc joins a boundary arc (tangentially)
- jumps down at times where two boundary arcs join together



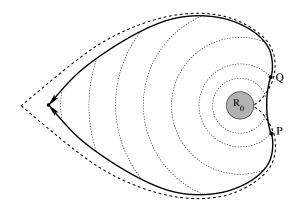
#### Junctions between different arcs



Two boundary arcs originating at the same point are not optimal



Non-parallel junctions between a free arc and a boundary arc are not optimal

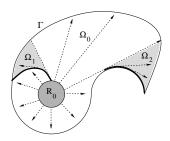


circle + two spirals

(is better than two spirals only)

#### Further classification

- blocking arcs  $\Gamma^b \doteq \Gamma \cap \partial R^{\Gamma}_{\infty}$
- delaying arcs  $\Gamma^d$

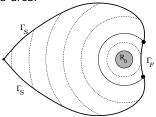


Necessary conditions for the optimality of delaying arcs (Tao Wang, 2010)

# Sufficient conditions for optimality ?

#### **Standard Isotropic Problem:**

- Fire starts on the unit disc, propagating with unit speed in all directions.
- Barrier can be constructed at speed  $\sigma > 2$ .
- Minimize the total burned area.

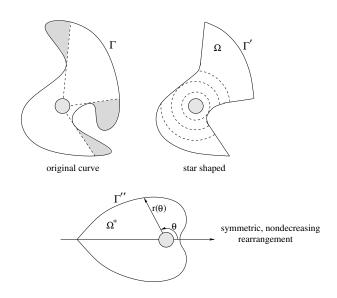


## **Theorem** (A.B. - T.Wang, 2010)

The barrier consisting of

#### circle + two logarithmic spirals

is optimal among all simple closed curves enclosing  $R_0$ 



## Symmetric rearrangement

$$\Omega = \left\{ (r\cos\theta, r\sin\theta); \ 0 \le r \le \rho(\theta) \right\}$$

$$\Omega^* = \left\{ (r\cos\theta, r\sin\theta); \ 0 \le r \le \tilde{\rho}^*(\theta) \right\},$$

 $\rho^*: [-\pi, \pi] \mapsto \mathbb{R}_+$  is the symmetric, nondecreasing rearrangement  $r(\cdot)$ .

$$\mathsf{meas}\Big(\{\theta\,;\ \ \rho^*(\theta)\ <\ c\}\Big)\ =\ \mathsf{meas}\Big(\{\theta\,;\ \ \rho(\theta)\ <\ c\}\Big)\qquad \mathsf{for\ all}\ c>0$$

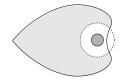
Polar coordinate representation:  $\theta \mapsto r(\theta)$  non-decreasing, for  $\theta \in [0, \pi]$ 

Admissibility constraint: 
$$m_1igg(\{x\in\Gamma\,;\;\;|x|\le 1+t\}igg)\le\sigma t$$

 $\mathsf{not} \ \mathsf{saturated} \ \Longrightarrow \mathsf{circumferences}$ 

saturated  $\implies$  logarithmic spirals





## **Numerical Simulations**

Assume: only blocking arcs, no delaying arcs.

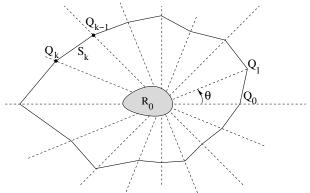
minimize total burned area:  $m_2\left(R_{\infty}^{\Gamma}\right)$ 

subject to

$$m_1\left(\Gamma \cup \overline{R^\Gamma(t)}\right) \leq \sigma t$$
 for all  $t \geq 0$ 

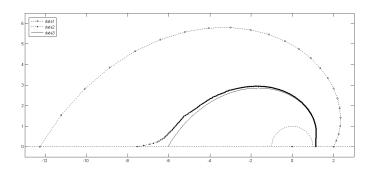
Approximate the barrier with a polygonal:

- fix an angle  $\theta = 2\pi/n$
- assign radii  $r_k = r(k\theta), \quad k = 1, \ldots, n$
- starting with an admissible polygonal, search for a local minimizer subject to admissibility constraints
- double number of nodes (replace n by 2n), repeat local minimization . . .



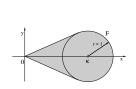
## 1. The isotropic case

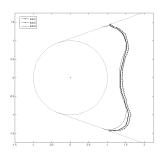
 $F(x) = R_0 = B_1$  (unit disc),  $\sigma = 4$ . Minimize: total burned area.



## 2. A non-isotropic case

$$F = \left\{ (\lambda x, \lambda y); (x-3)^2 + y^2 \le 1, \lambda \in [0,1] \right\}$$





Choose:  $\sigma = 4.1$ ,  $R_0 = \text{unit disc.}$ 

Analytic solution: A.Friend (2007). Numerical solution: T.Wang (2008)

## Some open problems

- 1 (Isotropic blocking problem). On the whole plane, assume:
- fire propagates with unit speed in all directions.

**Conjecture 1:** a blocking strategy exists if and only if the wall construction speed is  $\sigma > 2$ .

**2 (Sufficient conditions).** Not one single example is known where a blocking strategy can be proved to be optimal.

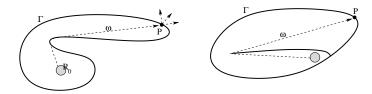
**Conjecture 2:** the "circle + two spirals" strategy is optimal for the isotropic problem.

Basic difficulty: delaying arcs

**3 (Existence of optimal strategies).** Determine whether an optimal strategy exists, in the general case where the velocity sets satisfy  $0 \in F(x)$  but without assuming  $B(0,\rho) \subset F(x)$  (so that fire propagation speed is not uniformly positive in all directions).

**4 (Regularity).** If the initial set  $R_0$  has a smooth boundary and the cost functions are smooth, what is the regularity of an optimal strategy? Does it produce a finite number of piecewise  $\mathcal{C}^1$  arcs? Is the optimal barrier connected? When is is useful to construct delaying arcs?

## On the minimal speed for a blocking strategy



If  $\Gamma$  is a simple closed curve, then the construction speed must be  $\sigma > 2$ .

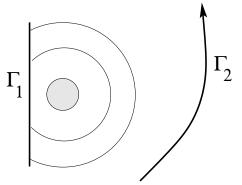
Indeed, let P= last brick of the wall. Then

$$m_1(\omega) < \frac{1}{2} m_1(\Gamma)$$

This estimate breaks down if  $\Gamma$  is not a simple closed curve.

## Alternative blocking strategies

Construct first a partial barrier, to slow down fire propagation. Then build a wall enclosing the fire.



(unlikely to succeed, with speed  $\sigma \leq 2$ )